

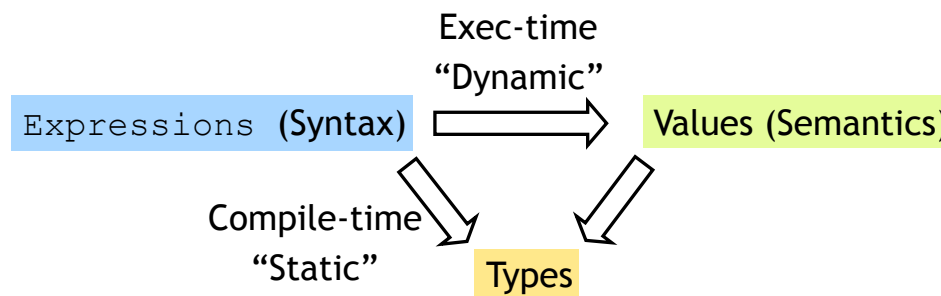
CSE 130 Programming Languages

Datatypes

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Recap: ML's Holy Trinity

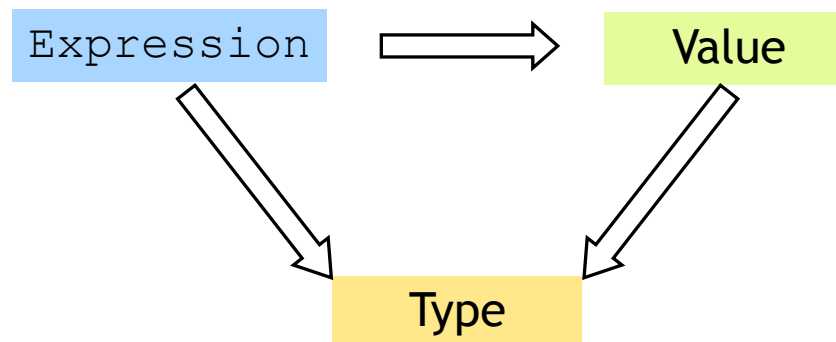


1. Programmer enters expression
2. ML checks if expression is "well-typed"
 - Using a precise set of rules, ML tries to find a unique type for the expression meaningful type for the expr
3. ML evaluates expression to compute value
 - Of the same "type" found in step 2

Story So Far...

- Simple Expressions
- Branches
- Let-Bindings ...
- Today:
 - Finish Crash Course
 - Datatypes

Next: functions, but remember ...



Everything is an expression
Everything has a value
Everything has a type

A function is a value!

A shorthand for function binding

```
# let neg = fun f -> fun x -> not (f x);
...
# let neg f x = not (f x);
val neg : int -> int -> bool = fn

# let is5gte = neg is5lt;
val is5gte : int -> bool = fn;
# is5gte 10;
val it : bool = false;
# is5gte 2;
val it : bool = true;
```

Put it together: a “filter” function

If arg “matches”
this pattern... ..then use
this Body Expr

```
- let rec filter f xs =
  match xs with
  | [] -> []
  | (x::xs') -> if f x
                 then x::(filter f xs')
                 else (filter f xs');;

val filter : ('a->bool)->'a list->'a list = fn

# let list1 = [1;31;12;4;7;2;10];;
# filter is5lt list1 ;;
val it : int list = [31;12;7;10]
# filter is5gte list1;;
val it : int list = [1;4;2]
# filter even list1;;
val it : int list = [12;4;2;10]
```

Put it together: a “partition” function

```
# let partition f l = (filter f l, filter (neg f) l);
val partition : ('a->bool)->'a list->'a list * 'a list =
fn

# let list1 = [1,31,12,4,7,2,10];
- ...
# partition is5lt list1 ;
val it : (int list * int list) = ([31,12,7,10],[1,2,10])

# partition even list1;
val it : (int list * int list) = ([12,4,2,10],[1,31,7])
```

A little trick ...

```
# 2 <= 3;; ...
val it : bool = true
# "ba" <= "ab";;
val it : bool = false

# let lt = (<) ;;
val it : 'a -> 'a -> bool = fn

# lt 2 3;;
val it : bool = true;
# lt "ba" "ab" ;;
val it : bool = false;

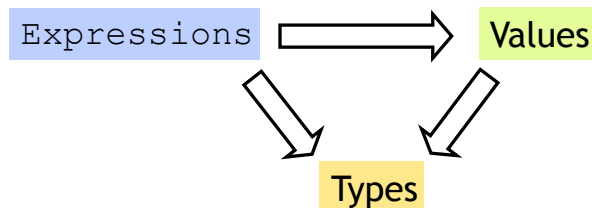
# let is5Lt = lt 5;
val is5lt : int -> bool = fn;
# is5lt 10;
val it : bool = true;
# is5lt 2;
val it : bool = false;
```

Put it together: a “quicksort” function

```
let rec sort xs =  
  match xs with  
  | [] -> []  
  | (h::t) -> let (l,r) = partition ((<) h) t in  
                (sort l)@(h::(sort r))
```

Now, lets begin at the beginning ...

What about more complex data ?



Many kinds of expressions:

1. Simple
2. Variables
3. Functions

News

- Ocaml-top issues?
- Please post questions to Piazza
- Seating: Don't Worry!

What about more complex data ?

- We've seen some **base** types and values:
 - Integers, Floats, Bool, String etc.
- Some ways to **build** up types:
 - Products (tuples), records, “lists”
 - Functions
- Design Principle: **Orthogonality**
 - Don't clutter **core language** with stuff
 - Few, powerful orthogonal building techniques
 - Put “**derived**” types, values, functions in **libraries**

What about more complex data ?

- We've seen some **base** types and values:
 - Integers, Floats, Bool, String etc.
- Some ways to **build** up types:
 - Products (tuples), records, "lists"
 - Functions

Next: Building datatypes

Three key ways to build complex types/values

1. "Each-of" types (**T1 * T2**)

Value of T contains value of T1 **and** a value of T2

2. "One-of" types

Value of T contains value of T1 **or** a value of T2

3. "Recursive"

Value of T contains (sub)-value of **same type** T

Next: Building datatypes

Three key ways to build complex types/values

1. "Each-of" types

Value of T contains value of T1 **and** a value of T2

2. "One-of" types

Value of T contains value of T1 **or** a value of T2

3. "Recursive"

Value of T contains (sub)-value of **same type** T

Suppose I wanted ...

... a program that processed lists of attributes

- Name (string)
- Age (integer)
- DOB (int-int-int)
- Address (string)
- Height (float)
- Alive (boolean)
- Phone (int-int)
- email (string)

Many kinds of attributes (too many to put in a record)

- can have multiple names, addresses, phones, emails etc.

Want to store them in a **list**. Can I ?

Suppose I wanted ...

Attributes:

- Name (string)
- Age (integer)
- DOB (int-int-int)
- Address (string)
- Height (real)
- Alive (boolean)
- Phone (int-int)
- email (string)

```
type attrib =  
  Name of string  
| Age of int  
| DOB of int*int*int  
| Address of string  
| Height of float  
| Alive of bool  
| Phone of int*int  
| Email of string;;
```

Constructing Datatypes

```
type t = C1 of t1 | C2 of t2 | ... | Cn of tn
```

t is a new datatype.

A value of type t is either:

- a value of type t_1 placed in a box labeled C_1
- Or a value of type t_2 placed in a box labeled C_2
- Or ...
- Or a value of type t_n placed in a box labeled C_n

Quiz: Here is a typedef ...

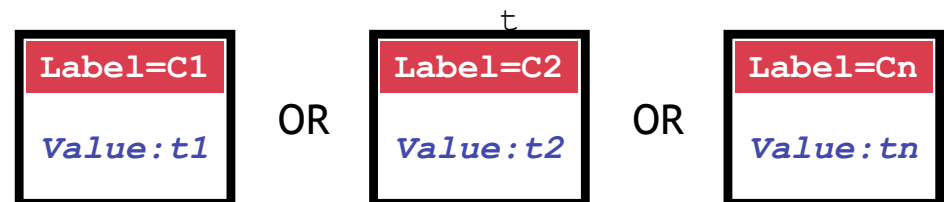
```
type attrib = Name of string  
             | Age of int  
             | Height of float
```

What is the type of: **Name “Tony Stark”**

- (a) Syntax Error
- (b) Type Error
- (c) **string**
- (d) **attrib**
- (e) 'a

Constructing Datatypes

```
type t = C1 of t1 | C2 of t2 | ... | Cn of tn
```



All have the type t

How to PUT values into box?



Question: Here is a typedef ...

```
type attrib = Name of string
             | Age of int
             | Height of float
```

What is the type of: **Age "Tony Stark"**

- (a) Syntax Error
- (b) Type Error
- (c) **string**
- (d) **attrib**
- (e) 'a

How to PUT values into box?

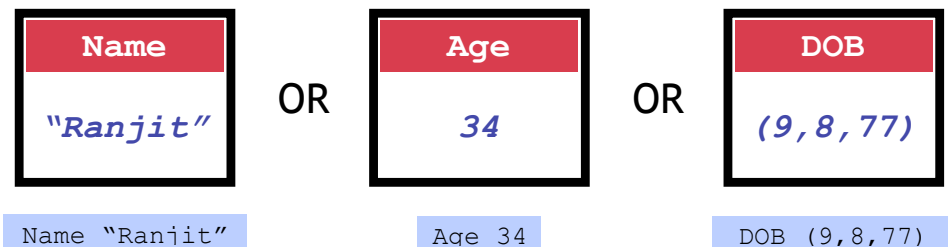
How to create values of type `attrib`?

```
# let a1 = Name "Ranjit";;
val x : attrib = Name "Ranjit"
# let a2 = Height 5.83;;
val a2 : attrib = Height 5.83
# let year = 1977 ;;
val year : int = 1977
# let a3 = DOB (9,8,year) ;;
val a3 : attrib = DOB (9,8,1977)
# let a_list = [a1;a2;a3];;
val a_list : attrib list = ...
```

```
type attrib =
  Name of string
  | Age of int
  | DOB of int*int*int
  | Address of string
  | Height of float
  | Alive of bool
  | Phone of int*int
  | Email of string;;
```

Constructing Datatypes

```
type attrib
= Name of string | Age of int | DOB of int*int*int
| Address of string | Height of float | Alive of bool
| Phone of int*int | Email of string;;
```



All have type **attrib**

One-of types

- We've defined a "one-of" type named `attrib`
- Elements are one of:
 - `string`,
 - `int`,
 - `int*int*int`,
 - `float`,
 - `bool` ...
- Can create uniform `attrib` lists
- Say I want a function to print attribs...

```
datatype attrib =  
  Name of string  
| Age of int  
| DOB of int*int*int  
| Address of string  
| Height of real  
| Alive of bool  
| Phone of int*int  
| Email of string;
```

Question: Here is a typedef ...

```
type attrib = Name of string  
             | Age of int  
             | Height of float
```

What is the type of:

```
[Name "Ranjit"; Age 35; Dob(9,8,77)]
```

- (a) Syntax Error
- (b) Type Error
- (c) `string * int * (int*int*int) list`
- (d) 'a list
- (e) `attrib list`

How to TEST & TAKE whats in box?



Is it a ...
string?
or an
int?
or an
int*int*int?
or ...

How to TEST & TAKE whats in box?



Look at TAG!

Question: Here is a typedef ...

```
type attrib = Name of string | Age of int | ...
```

What does this evaluate to?

```
let welcome a = match a with
    | Name s -> s
in welcome (Name "Ranjit")
```

- (a) Name "Ranjit" : 'a
- (b) Type Error
- (c) Name "Ranjit" : attrib
- (d) "Ranjit" : string
- (e) Runtime Error

How to tell whats in the box ?

```
type attrib =
  Name of string
| Age of int
| DOB of int*int*int
| Address of string
| Height of float
| Alive of bool
| Phone of int*int

match e with
| Name s -> ...(*s: string *)
| Age i -> ...(*i: int *)
| DOB(d,m,y)-> ...(*d: int,m: int,y: int*)
| Address a -> ...(*a: string*)
| Height h -> ...(*h: int *)
| Alive b -> ...(*b: bool*)
| Phone(y,r)-> ...(*a: int, r: int*)
```

Pattern-match expression: check if e is of the form ...

- On match:
 - value in box bound to pattern variable
 - matching result expression is evaluated
- Simultaneously test and extract contents of box

How to tell whats in the box ?

```
match e with
| Name s -> printf "%s" s
| Age i -> printf "%d" i
| DOB(d,m,y) -> printf "%d/%d/%d" d m y
| Address s -> printf "%s" s
| Height h -> printf "%f" h
| Alive b -> printf "%b" b s
| Phone(a,r) -> printf "(%d)-%d" a r
```

Pattern-match expression: check if e is of the form ...

- On match:
 - value in box bound to pattern variable
 - matching result expression is evaluated
- Simultaneously test and extract contents of box

Question: Here is a typedef ...

```
type attrib = Name of string | Age of int | ...
```

What does this evaluate to?

```
let welcome a = match a with
    | Name s -> s
in welcome (Age 34)
```

- (a) Name "Ranjit" : 'a
- (b) Type Error
- (c) Name "Ranjit" : attrib
- (d) "Ranjit" : string
- (e) Runtime Error

How to tell whats in the box

```
# match (Name "Ranjit") with
| Name s -> printf "Hello %s\n" s
| Age i  -> printf "%d years old" i
;;
```

```
Hello Ranjit
- : unit = ()
```

First case matches the tag (Name)

Evals branch with **s** “bound” to string contents

How to TEST & TAKE whats in box?



BEWARE!!
Be sure to
handle all
TAGS!

Beware! Handle All TAGS!

```
# match (Name "Ranjit") with
| Age i  -> Printf.printf "%d" i
| Email s -> Printf.printf "%s" s
;;
```

Exception: Match Failure!!

None of the cases matched the tag (Name)

Causes nasty **Run-Time Error**

Compiler To The Rescue!!

```
# let printAttrib a =
  match a with
  | Name s -> Printf.printf "%s" s
  | Age i  -> Printf.printf "%d" i
  | DOB (d,m,y) -> Printf.printf "%d / %d / %d" d m y
  | Address addr -> Printf.printf "%s" addr
  | Height h -> Printf.printf "%f" h
  | Alive b -> Printf.printf "%b" b
  | Email e -> Printf.printf "%s" e
  ;;
Warning P: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
Phone (_, _)
```

Compile-time checks for:

missed cases: ML warns if you miss a case!

Q: What does this evaluate to ?

```
type attrib = Name of string | ...
let welcome a = match a with
  | Name s -> "Hello!" ^ s
  | Name s -> "Welcome!" ^ s
in welcome (Name "Mickey")
```

- (a) Type Error
- (b) "Welcome!Mickey" : string
- (c) Runtime Error
- (d) "Hello!Mickey" : string
- (e) "Hello!MickeyWelcome!Mickey"Ranjit" : string

Benefits of match-with

```
match e with
  C1 x1 -> e1
| C2 x2 -> e2
| ...
| Cn xn -> en
```

```
type t =
  C1 of t1
| C2 of t2
| ...
| Cn of tn
```

1. Simultaneous test-extract-bind
2. Compile-time checks for:
 - missed cases: ML warns if you miss a `t` value
 - redundant cases: ML warns if a case never matches

Compiler To The Rescue!!

```
# let printAttrib a =
  match a with
  | Name s -> Printf.printf "%s" s
  | Age i -> Printf.printf "%d" i
  | DOB (d,m,y) -> Printf.printf "%d / %d / %d" d m y
  ...
  | Age i -> Printf.printf "%d" i
;;
```

Warning U: this match case is unused.

Compile-time checks for:

redundant cases: ML warns if a case never matches

match-with is an Expression

```
match e with
  C1 x1 -> e1
| C2 x2 -> e2
| ...
| Cn xn -> en
```

Q: What does this evaluate to ?

```
type attrib = Name of string | Age of int | ...
```

```
let welcome a = match a with  
    | Name s -> s  
    | Age i  -> i  
in welcome (Name "Ranjit")
```

- (a) "Ranjit" : string
- (b) Type Error
- (c) Name "Ranjit" : attrib
- (d) Runtime Error

Next: Building datatypes

Three key ways to build complex types/values

1. "Each-of" types $t1 * t2$

Value of T contains value of T1 **and** a value of T2

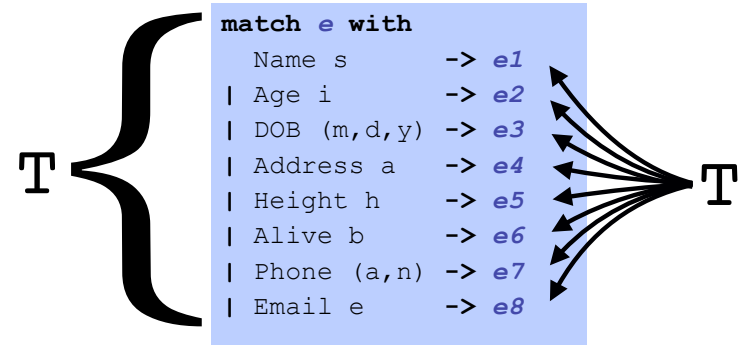
2. "One-of" types $\text{type } t = C1 \text{ of } t1 \mid C2 \text{ of } t2$

Value of T contains value of T1 **or** a value of T2

3. "Recursive" type

Value of T contains (sub)-value of **same type** T

match-with is an Expression



Type Rule

- $e1, e2, \dots, en$ must have same type T
- Type of whole expression is T

"Recursive" types

```
type nat = Zero | Succ of nat
```

“Recursive” types

```
type nat = Zero | Succ of nat
```

Wait a minute! **Zero** of what ?!

“Recursive” types

```
type nat = Zero | Succ of nat
```

Wait a minute! **Zero** of what ?!

Relax.

Means “empty box with label **Zero**”

“Recursive” types

```
type nat = Zero | Succ of nat
```

What are values of `nat` ?

“Recursive” types

```
type nat = Zero | Succ of nat
```

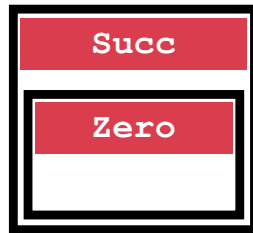
What are values of `nat` ?



“Recursive” types

```
type nat = Zero | Succ of nat
```

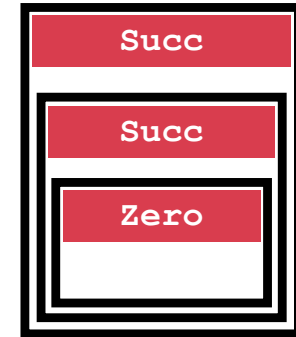
What are values of `nat` ?
One `nat` contains another!



“Recursive” types

```
type nat = Zero | Succ of nat
```

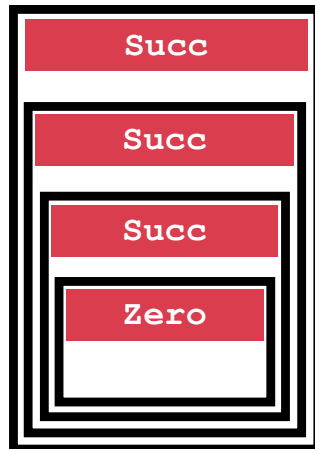
What are values of `nat` ?
One `nat` contains another!



“Recursive” types

```
type nat = Zero | Succ of nat
```

What are values of `nat` ?
One `nat` contains another!

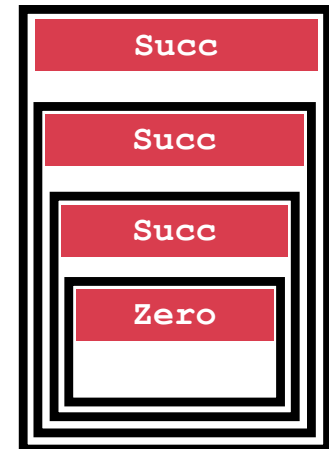


“Recursive” types

```
type nat = Zero | Succ of nat
```

What are values of `nat` ?
One `nat` contains another!

`nat` = recursive type



Next: Building datatypes

Three key ways to build complex types/values

1. “Each-of” types **t1 * t2**

Value of T contains value of T1 **and** a value of T2

2. “One-of” types **type t = C1 of t1 | C2 of t2**

Value of T contains value of T1 **or** a value of T2

3. “Recursive” type **type t = ... | C of (...*t)**

Value of T contains (sub)-value of **same type T**

Next: Lets get cosy with Recursion

Code Structure = Type Structure!!!

Next: Lets get cosy with Recursion

Recursive Code Mirrors Recursive Data

`to_int : nat -> int`

```
type nat =  
| Zero  
| Succ of nat
```

```
let rec to_int n =
```

to_int : nat -> int

```
type nat =  
  Base pattern | Zero  
  Inductive pattern | Succ of nat
```

```
let rec to_int n =
```

to_int : nat -> int

```
type nat =  
  Base pattern | Zero  
  Inductive pattern | Succ of nat
```

```
let rec to_int n = match n with  
  Base pattern | Zero ->  
  Inductive pattern | Succ m ->
```

to_int : nat -> int

```
type nat =  
  Base pattern | Zero  
  Inductive pattern | Succ of nat
```

```
let rec to_int n = match n with  
  Base pattern | Zero -> 0 Base Expression  
  Inductive pattern | Succ m -> 1 + to_int m Inductive Expression
```

Q: What does this evaluate to ?

```
let rec foo n =  
  if n <= 0 then Zero else Succ(foo(n-1))  
in foo 2
```

- (a) Zero : nat
- (b) Type Error
- (c) 2 : nat
- (c) Succ(Zero) : nat
- (c) Succ(Succ(Zero)) : nat

of_int : int -> nat

```
type nat =  
| Zero  
| Succ of nat
```

```
let rec of_int n =
```

of_int : int -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec of_int n =
```

of_int : int -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec of_int n =  
  if n <= 0 then  
  else
```

Base pattern if n <= 0 then
Inductive pattern else

of_int : int -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec of_int n =  
  if n <= 0 then  
    Zero Base Expression  
  else  
    Succ (of_int (n-1)) Inductive Expression
```

Base pattern if n <= 0 then
Inductive pattern else

plus : nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

```
let rec plus (n,m) =
```

plus : nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec plus (n,m) =
```

plus : nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec plus (n,m) =  
  match m with  
  | Zero ->  
  | Succ m' ->
```

Base pattern | Zero ->
Inductive pattern | Succ m' ->

plus : nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec plus (n,m) =  
  match m with  
  | Zero -> n  
  | Succ m' -> Succ (plus (n,m'))
```

Base pattern | Zero -> n *Base Expression*
Inductive pattern | Succ m' -> Succ (plus (n,m')) *Inductive Expression*

times: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

```
let rec times (n,m) =
```

times: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec times (n,m) =
```

times: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec times (n,m) =  
  match m with  
  | Zero ->  
  | Succ m' ->
```

Base pattern | Zero ->
Inductive pattern | Succ m' ->

plus : nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec times (n,m) =  
  match m with  
  | Zero -> Zero  
  | Succ m' -> plus n (times (n,m'))
```

Base pattern | Zero -> Zero
Inductive pattern | Succ m' -> plus n (times (n,m'))

Base Expression
Inductive Expression

minus: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

```
let rec minus (n,m) =
```

times: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec minus (n,m) =
```

times: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec minus (n,m) =  
  match (n, m) with  
  | (_, Zero) -> n  
  | (Succ n', Succ m') -> minus(n', m')
```

Base pattern | (_, Zero) -> n
Inductive pattern | (Succ n', Succ m') -> minus(n', m')

times: nat*nat -> nat

```
type nat =  
| Zero  
| Succ of nat
```

Base pattern | Zero
Inductive pattern | Succ of nat

```
let rec minus (n,m) =  
  match (n, m) with  
  | (_, Zero) -> n  
  | (Succ n', Succ m') -> minus(n', m')
```

Base pattern | (_, Zero) -> n *Base Expression*
Inductive pattern | (Succ n', Succ m') -> minus(n', m') *Inductive Expression*

Lists are recursive types!

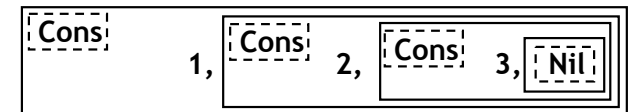
Next: Lets get cosy with Recursion

Recursive Code Mirrors Recursive Data

```
type int_list =
  Nil
| Cons of int * int_list
```

Think about this! What are values of `int_list` ?

`Cons(1,Cons(2,Cons(3,Nil)))` `Cons(2,Cons(3,Nil))` `Cons(3,Nil)` `Nil`



Lists aren't built-in !

```
datatype int_list =
  Nil
| Cons of int * int_list
```

Lists are a **derived** type: built using elegant core!

1. Each-of
2. One-of
3. Recursive

`::` is just a pretty way to say “Cons”

`[]` is just a pretty way to say “Nil”

Some functions on Lists : Length

```
let rec len l =
  match l with
  Nil -> 0
| Cons(h,t) -> 1 + (len t)
```

Base pattern (points to Nil) *Base Expression* (points to 0)

Ind pattern (points to Cons(h,t)) *Inductive Expression* (points to 1 + (len t))

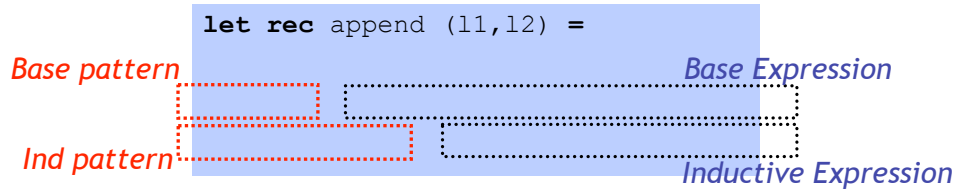
```
let rec len l =
  match l with
  Nil -> 0
| Cons(_,t) -> 1 + (len t)
```

Matches everything, no binding

```
let rec len l =
  match l with
  Cons(_,t) -> 1 + (len t)
| _ -> 0
```

Pattern-matching in order
- Must match with Nil

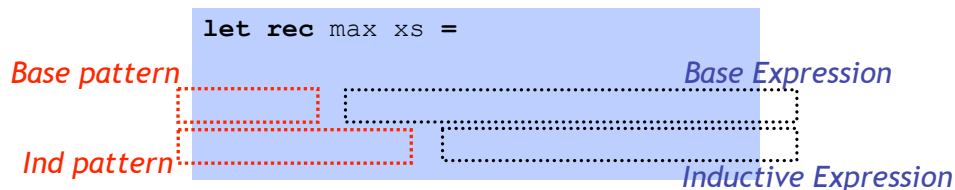
Some functions on Lists : Append



- Find the right induction strategy
 - Base case: pattern + expression
 - Induction case: pattern + expression

Well designed datatype gives strategy

Some functions on Lists : Max



- Find the right induction strategy
 - Base case: pattern + expression
 - Induction case: pattern + expression

Well designed datatype gives strategy

null, hd, tl are all functions ...

Bad ML style: More than aesthetics !

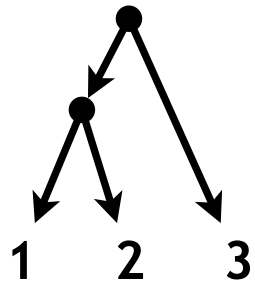
Pattern-matching better than test-extract:

- ML checks all cases covered
- ML checks no redundant cases
- ...at compile-time:
 - fewer errors (crashes) during execution
 - get the bugs out ASAP!

Next: Lets get cosy with Recursion

Recursive Code Mirrors Recursive Data

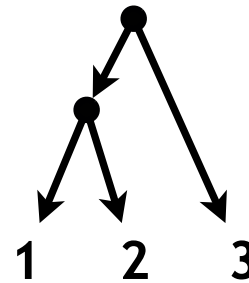
Q: How is this tree represented ?



```
type tree =  
| Leaf of int  
| Node of tree*tree
```

- (a) (1, 2), 3
- (b) (Leaf 1, Leaf 2), Leaf 3
- (c) Node (Node (Leaf 1, Leaf 2), Leaf 3)
- (d) Node ((Leaf 1, Leaf 2), Leaf 3)
- (e) None of the above

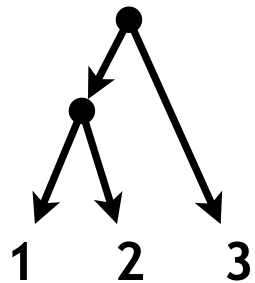
Representing Trees



```
type tree =  
| Leaf of int  
| Node of tree*tree
```

Leaf 1

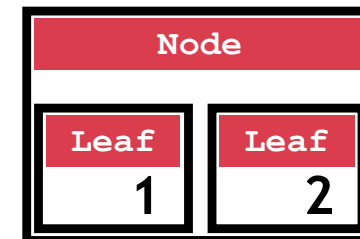
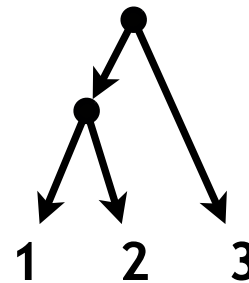
Representing Trees



Leaf 2

```
type tree =  
| Leaf of int  
| Node of tree*tree
```

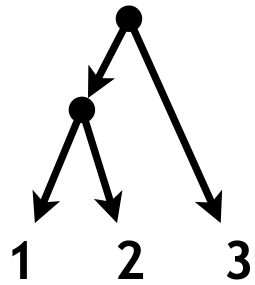
Representing Trees



Node(Leaf 1, Leaf 2)

```
type tree =  
| Leaf of int  
| Node of tree*tree
```

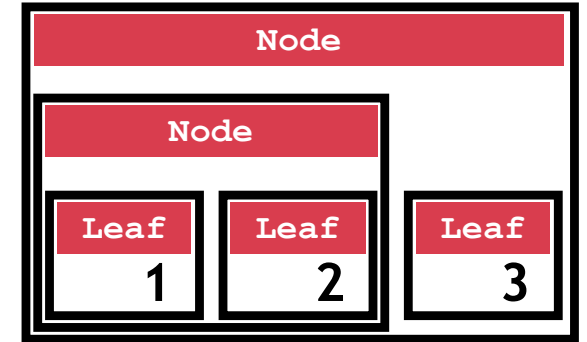
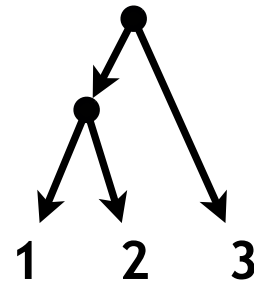
Representing Trees



```
type tree =  
| Leaf of int  
| Node of tree*tree
```

Leaf 3

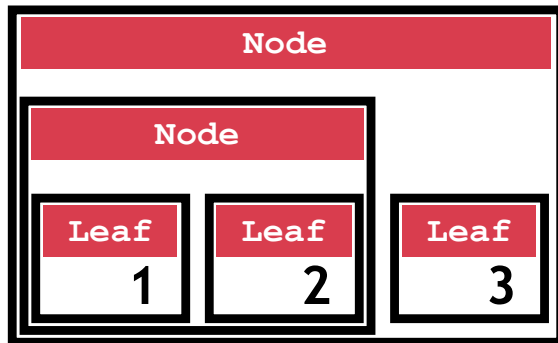
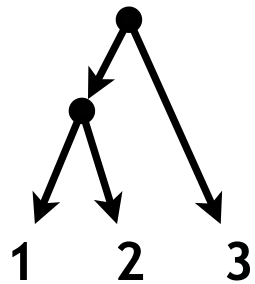
Representing Trees



```
type tree =  
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| Node of tree*tree
```

Node(Node(Leaf 1, Leaf 2), Leaf 3)

Representing Trees



```
type tree =  
| Leaf of int  
| Node of tree*tree
```

Node(Node(Leaf 1, Leaf 2), Leaf 3)

Next: Lets get cosy with Recursion

Recursive Code Mirrors Recursive Data

Q: What does this evaluate to ?

```
let rec foo t = match t with
| Leaf n      -> 1
| Node (t1, t2) -> foo t1 + foo t2
```

```
foo (Node(Node(Leaf 1,Leaf 2),Leaf 3))
```

- (a) **Type Error**
- (b) **1 : int**
- (c) **3 : int**
- (d) **6 : int**

sum_leaf: tree -> int

“Sum up the leaf values”. E.g.

```
# let t0 = Node(Node(Leaf 1, Leaf 2), Leaf 3);;
# sum_leaf t0 ;;
- : int = 6
```

sum_leaf: tree -> int

```
type tree =
| Leaf of int
| Node of tree*tree
```

```
let rec sum_leaf t =
```

sum_leaf: tree -> int

```
type tree =
| Leaf of int
| Node of tree*tree
```

Base pattern | Leaf of int
Inductive pattern | Node of tree*tree

```
let rec sum_leaf t =
```


sum_leaf: tree -> int

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Base pattern | Leaf of int
Inductive pattern | Node of tree*tree

```
let rec sum_leaf t =  
  match t with  
  | Leaf n ->  
  | Node(t1,t2)->
```

Base pattern | Leaf n ->
Inductive pattern | Node(t1,t2)->

sum_leaf: tree -> int

```
type tree =  
  | Leaf of int  
  | Node of tree*tree
```

Base pattern | Leaf of int
Inductive pattern | Node of tree*tree

```
let rec sum_leaf t =  
  match t with  
  | Leaf n -> n Base Expression  
  | Node(t1,t2)-> sum_leaf t1 + sum_leaf t2  
  Inductive Expression
```

Base pattern | Leaf n -> n *Base Expression*
Inductive pattern | Node(t1,t2)-> sum_leaf t1 + sum_leaf t2
Inductive Expression

Recursive Code Mirrors Recursive Data

Code almost writes itself!

Another Example: Calculator

Want an arithmetic calculator to evaluate expressions like:

- 4.0 + 2.9
- 3.78 - 5.92
- (4.0 + 2.9) * (3.78 - 5.92)

Another Example: Calculator

Want an arithmetic calculator to evaluate expressions like:

- $4.0 + 2.9 \implies 6.9$
- $3.78 - 5.92 \implies -2.14$
- $(4.0 + 2.9) * (3.78 - 5.92) \implies -14.766$

Whats a ML **TYPE** for **REPRESENTING** expressions ?

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Whats a ML **FUNCTION** for **EVALUATING** expressions ?

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type expr =  
| Num of float  
| Add of expr*expr  
| Sub of expr*expr  
| Mul of expr*expr
```

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```
type expr =
| Num of float
| Add of expr*expr
| Sub of expr*expr
| Mul of expr*expr

let rec eval e = match e with
| Num f      ->
| Add (e1,e2)->
| Sub (e1,e2)->
| Mul (e1,e2)->
```

Another Example: Calculator

Want an arithmetic calculator to evaluate expressions like:

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Whats a ML **FUNCTION** for **EVALUATING** expressions ?

```
type expr =
| Num of float
| Add of expr*expr
| Sub of expr*expr
| Mul of expr*expr

let rec eval e = match e with
| Num f      -> f
| Add (e1,e2)-> eval e1 +. eval e2
| Sub (e1,e2)-> eval e1 -. eval e2
| Mul (e1,e2)-> eval e1 *. eval e2
```

